

METHOD AND APPARATUS FOR THE USE OF MICRO-TUNNELS IN A COMMUNICATIONS SYSTEM

Claim of Priority under 35 U.S.C. §119

[0001] The present Application for Patent claims priority to Provisional Application No. 60/418,815, entitled, "Micro-Tunnels," filed 15 October 2002, and assigned to the assignee hereof and hereby expressly incorporated by reference herein.

BACKGROUND

Field

[0002] The present invention relates generally to data packet transmission and specifically to the use of micro-tunnels.

Background

[0003] Internet Protocol (IP) "tunnels" have become a widespread mechanism to transport data units, referred to as datagrams, over the Internet. Using Tunneling involves incorporating an original IP packet inside of another IP packet. Tunneling also has additional connotations about changing the effects of Internet routing on the original IP packet.

[0004] Typically, a tunnel is used to augment and modify the behavior of the deployed routing architecture, such as in multicast routing, mobile IP, and Virtual Private Network (VPN). From the perspective of traditional best-effort IP packet delivery, a tunnel behaves as any other link. Packets enter one end of the tunnel, and are delivered to the other end unless resource overload or error causes them to be lost.

[0005] Information may be encapsulated and routed through a tunnel. In the most general case, a system has a packet, which is referred to as a payload packet, which needs to be encapsulated and routed. The payload packet is first encapsulated in a Generic Routing Encapsulation (GRE) packet, which possibly also includes a routing. The resulting GRE packet may then be encapsulated in some other protocol and then forwarded. This outer protocol is referred to as the delivery protocol.

[0006] For mobile IP, a wireless system interfaces with an IP network. Tunnels are used for transporting data from the IP network to infrastructure elements in the wireless

system. The data may involve multiple streams of data for transmission to and/or from a same mobile node. In this case, the system must establish individual tunnels for each stream.

[0007] In mobile IP the home agent associated with the mobile node redirects packets from the home network to the care-of address by constructing a new IP header containing the mobile node's care-of address as the destination IP address. The home agent is a router on a mobile node's home network maintaining information about the device's current location, as identified in its care-of address. The care-of address is a temporary IP address for a mobile node enabling message delivery when the device is connecting from somewhere other than its home network. The care-of address identifies a mobile node's current point of attachment to the Internet and makes it possible to connect from a different location without changing the device's home address (permanent IP address). The new header then shields or encapsulates the original packet, causing the mobile node's home address to have no effect on the encapsulated packet's routing until it arrives at the care-of address. Such encapsulation is also called tunneling, which suggests the packet burrows through the Internet, bypassing the usual effects of IP routing.

[0008] In a mobile IP environment, there is a need to identify multiple tunnels each associated with a same mobile node. Further, there is a need for flexible tunnel set up which optimizes the resources of the system.

BRIEF DESCRIPTION OF THE DRAWINGS

[0009] FIG. 1 is a wireless communication system supporting mobile IP.

[0010] FIG. 2 is a Generic Routing Encapsulation (GRE) format.

[0011] FIG. 3 is a GRE key format.

[0012] FIG. 4 is an illustration of the GRE key space.

DETAILED DESCRIPTION

[0013] The field of wireless communications has many applications including, e.g., cordless telephones, paging, wireless local loops, Personal Digital Assistants (PDAs), Internet telephony, and satellite communication systems. A particularly important application is cellular telephone systems for mobile subscribers. As used herein, the term "cellular" system encompasses both cellular and Personal Communication

Services (PCS) frequencies. Various over-the-air interfaces have been developed for such cellular telephone systems including, e.g., Frequency Division Multiple Access (FDMA), Time Division Multiple Access (TDMA), and Code Division Multiple Access (CDMA). In connection therewith, various domestic and international standards have been established including, e.g., Advanced Mobile Phone Service (AMPS), Global System for Mobile (GSM), and Interim Standard 95 (IS-95). IS-95 and its derivatives, IS-95A, IS-95B, ANSI J-STD-008 (often referred to collectively herein as IS-95), and proposed high-data-rate systems are promulgated by the Telecommunication Industry Association (TIA) and other well known standards bodies.

[0014] Cellular telephone systems configured in accordance with the use of the IS-95 standard employ CDMA signal processing techniques to provide highly efficient and robust cellular telephone service. An exemplary system utilizing CDMA techniques is the cdma2000 ITU-R Radio Transmission Technology (RTT) Candidate Submission (referred to herein as cdma2000), issued by the TIA. The standard for cdma2000 is given in the draft versions of IS-2000 and has been approved by the TIA and 3GPP2. Another CDMA standard is the W-CDMA standard, as embodied in 3rd Generation Partnership Project "3GPP," Document Nos. 3G TS 25.211, 3G TS 25.212, 3G TS 25.213, and 3G TS 25.214.

[0015] A cellular communication system supporting mobile IP is illustrated in FIG. 1. System 100 supports communications of packets of data, wherein a packet is a logical grouping of information including a header containing control information and user data. Packets most often are used to refer to network layer units of data. Note the terms datagram, frame, message, and segment also are used to describe logical information groupings at various layers of the Open Systems Interconnection (OSI) reference model.

[0016] In the system 100, a Packet Data Service Node (PDSN) 102 interfaces between the wireless communication system and an IP network. In the mobile IP environment, the PDSN may also be referred to as a Foreign Agent (FA). The Home Agent (HA) is the node in the home network of the mobile node effectively causing the mobile node to be reachable at a home address even when the mobile node is not attached to the home network.

[0017] Continuing with FIG. 1, the PDSN 102 communicates with the various Mobile Nodes (MNs) via a Packet Control Function/Base Station Controller (PCF/BSC) 104. Note the PCF and BSC may reside in separate infrastructure elements or may be

combined in one element as illustrated in FIG. 1. The PDSN 102 provides communications for MN 108, 110 via the PCF/BSC 104. A Mobile Switching Center (MSC) is also in communication with PCF/BSC 104.

[0018] For a typical packet data communication, the PCF/BSC 104 sends an A11-Registration Request message to the PDSN 102 to establish an A10/11 interface between itself and the PDSN. The various interfaces refer to the communication links or sessions between the infrastructure elements. The A11 interface is generally identified as the link between the PDSN 102 and the PCF/BSC 104.

[0019] The PCF/BSC 104 binds the mobile station identifier, e.g. a Mobile Identification Number (MIN) such as an International Mobile Subscriber Identity (IMSI), to a Packet Session Identifier (PSI) unique within the PCF/BSC 104. The IMSI is a number used to uniquely identify personal mobile stations (i.e., mobile nodes). In some cases, ambiguity might arise when using only the 10-digit MIN. In one system, the first three (most significant) decimal numbers of the IMSI are the Mobile Country Code (MCC); the remaining digits are the National Mobile Station Identity (NMSI).

[0020] For each data communication to MN 108 or MN 110, the PCF/BSC 104 establishes a separate link. When the PCF/BSC 104 establishes the link, the PCF/BSC 104 includes the PSI in the A11-Registration-Request message that is sent to the PDSN. In this way, a communication intended for a given MN, such as MN 108 or MN 110, is processed via the designated link. As the number of data services increases, a MN may desire to have multiple data communications concurrently. In this case, the PCF/BSC 104 seeks to establish a link for each communication.

[0021] As described herein “micro-tunnels” are logical connections between the PDSN 102 and the PCF/BSC 104 that are identified by a source IP address and a destination IP address. For example, the source IP address may be identified as “src_ip_address,” and the destination IP address may be identified as “dest_ip_address.” A micro-tunnel is then designated by the following:

<src_ip_address = PDSN_IP, dest_ip_address = PCF_IP, GRE_key>.

In this context, the source refers to the PDSN 102, the destination refers to the PCF/BSC 104. Note, micro-tunnels are independent of the air-interface service instances (i.e., no one-to-one mapping is assumed between the micro-tunnels and the air-interface service instances).

[0022] Each micro-tunnel is assigned a separate communication for a given MN. As illustrated in FIG. 1, multiple micro-tunnels may be established for one MN. In the example of FIG. 1, three micro-tunnels are established for three separate communications to MN 108, while two micro-tunnels are established for two separate communications to MN 110. However, a single communication may utilize one or more micro-tunnels. The micro-tunnels are established when the PCF/BSC 104 sends a message to the PDSN 102. Specifically, the message is an A11-Registration-Request.

[0023] Once a micro-tunnel is established, a communication may be transmitted via the established micro-tunnel. There is not necessarily a one-to-one mapping between the air-interface service instances and micro-tunnels.

[0024] The micro-tunnel serves the following purposes:

- Identify the PPP context;
- Identify the IP context; and
- Differentiate services.

The following discussion details each of these micro-tunnel functions.

PPP Context:

[0025] The micro-tunnels are used by the PDSN 102 to indicate to the Radio Access Network (RAN) 120 whether the data packets carried by the micro-tunnel may be dropped or not. In lieu of such indication, the RAN 120 may decide to drop packets if they get too stale. For example, if stateful compression or encryption is used, dropping packets in the RAN 120 may cause problems for de-compression. State is a collection of information maintained by an entity. Stateful encryption or compression means the encryptor/decryptor or the compressor/decompressor maintains state information. Dropped packets therefore will impact the compressor/decompressor processes. In such a case, the PDSN 102 selects a "do not drop" attribute for the micro-tunnel.

[0026] The micro-tunnels are further used by the PDSN 102 to indicate those packets transported via a given micro-tunnel are to be treated differently from the other packets transported via another micro-tunnel. For example, the PDSN 102 may indicate packets carried by a first micro-tunnel have a particular header compression, such as zero-byte-header compression. The PCF/BSC 104 then interprets this information and uses a Radio Link Protocol (RLP)-free service instance to carry these packets.

IP Context:

- [0027] Using micro-tunnels the PDSN 102 indicates to the RAN 120 that re-ordering of data packets is allowed across micro-tunnels but not within the micro-tunnels. This approach is consistent with recommendations in section 4.1 of "Differentiated Services and Tunnels" by D. Black, published October 2000, and identified as RFC 2983 by the Internet Engineering Task Force (IETF). In some situations, it may be desirable to enable reordering among packets in different micro-tunnels to coexist with an absence of packet reordering within each individual micro-tunnel.
- [0028] In a first scenario, a node supporting various quality of service requirements and discriminate among packets, such as a Differentiated Services (DS) node, is instructed not to re-order packets belonging to the same micro-flow and the same quality of service requirements, such as an Assured Forwarding (AF) class. Note, DS and AF classes are defined in: (1) "Assured Forwarding PHB Group" by J. Heinanen et al., published June 1999 and identified as RFC 2597; (2) "Definition of the Differentiated Services Field (DS Field) in the IPv4 and IPv6 Headers" by K. Nichols, published December 1998, and identified as RFC 2474; and (3) "An Architecture for Differentiated Services" by S. Blake, published December 1998, and identified as RFC 2475. Each of the RFC documents referenced herein is provided by the Network Working Group of the Internet Engineering Task Force (IETF).
- [0029] Differentiated Services (DS) are intended to provide a framework and building blocks to enable deployment of scalable service discrimination in the Internet. The differentiated services approach aims to speed deployment by separating the architecture into two major components, one of which is fairly well-understood and the other of which is just beginning to be understood. Packet forwarding is a task which is performed on a per-packet basis as quickly as possible. Forwarding uses the packet header to find an entry in a routing table specifying the packet's output interface. Routing sets the entries in the table and may need to reflect a range of transit and other policies as well as to keep track of route failures. Routing tables are maintained as a background process to the forwarding task. A Differentiated Services Domain is a contiguous portion of the Internet over which a consistent set of differentiated services policies are administered in a coordinated fashion. A differentiated services domain may represent different administrative domains or autonomous systems, different trust

regions, different network technologies (e.g., cell/frame), hosts and routers, etc. Alternate embodiments may apply alternate methods whereby packets are discriminated among and unique treatment applied thereto. Alternate services also provide quality of service variations to different data packets.

[0030] Assured Forwarding of IP packets over the Internet is desirable, for example, when a company uses the Internet to interconnect to geographically distributed sites and wants an assurance that IP packets within this intranet are forwarded with high probability. In this situation, it is desirable for the network to not reorder packets belonging to the same microflow, wherein a microflow: is a single instance of an application-to-application flow of packets which is identified by source address, destination address, protocol id, and source port, destination port (where applicable).

[0031] Assured Forwarding (AF) grouping provides a means for a provider DS domain to offer different levels of forwarding assurances for IP packets received from a customer DS domain. Four AF classes are defined, wherein each AF class is, in each DS node, allocated a certain amount of forwarding resources, such as buffer space and bandwidth. IP packets wishing to use the services provided by the AF group are assigned by the customer or the provider DS domain into one or more of these AF classes according to the services to which the customer has subscribed.

[0032] Within each AF class, IP packets are marked with one of three possible drop precedence values. In case of congestion, the drop precedence of a packet determines the relative importance of the packet within the AF class. A congested DS node tries to protect packets with a lower drop precedence value from being lost by preferably discarding packets with a higher drop precedence value.

[0033] In a DS node, the level of forwarding assurance of an IP packet thus depends on: (1) the amount of forwarding resources allocated to the AF class to which the packet belongs, (2) the current load of the AF class, and, in case of congestion within the class, (3) the drop precedence of the packet.

[0034] For example, if traffic conditioning actions at the ingress of the provider DS domain make sure an AF class in the DS nodes is only moderately loaded by packets with the lowest drop precedence value and is not overloaded by packets with the two lowest drop precedence values, then the AF class may offer a high level of forwarding assurance.

- [0035] In another embodiment, the Assured Forwarding (AF) group provides forwarding of IP packets in N independent AF classes. Within each AF class, an IP packet is assigned one of M different levels of drop precedence. An IP packet belonging to an AF class i and has drop precedence j is marked with the AF codepoint AFij, where $1 \leq i \leq N$ and $1 \leq j \leq M$. Currently, four classes (N=4) with three levels of drop precedence in each class (M=3) are defined for general use. More AF classes or levels of drop precedence may be defined for local use.
- [0036] The identity of the micro-flow is hidden (due to GRE encapsulation) on the R-IP interface between the RAN and the PDSN. Therefore, the DS nodes between the PDSN and RAN cannot distinguish different micro-flows from each other unless the PDSN uses a micro-tunnel for each micro-flow in order to satisfy the in-sequence delivery requirement. Another example of the flows which are sensitive to re-ordering is flows protected by IPsec.
- [0037] The GRE format is illustrated in FIG. 2. The data packet format includes a delivery header 202, a GRE header 204 and a payload packet 206. The GRE header 204 may include a key field containing a four octet number which was inserted by the encapsulator. The key may be used by the receiver to authenticate the source of the packet. In one embodiment, the key field is made up of two fields. Also, the GRE header 204 may include a sequence number field. The sequence number field contains an unsigned 32 bit integer which is inserted by an encapsulator. It may be used by the receiver to establish the order in which packets have been transmitted from the encapsulator to the receiver.
- [0038] In another scenario, certain packets, such as Layer 2 Tunneling Protocol (L2TP) packets and IPsec packets, should not be re-ordered. By using a separate micro-tunnel for these types of traffic, the PDSN 102 instructs the PCF/BSC 104 that re-ordering is allowed among the IPsec/L2TP traffic, but not within a micro-tunnel. Note, L2TP is an industry-standard Internet tunneling protocol. Unlike Point-to-Point Tunneling Protocol (PPTP), L2TP does not require IP connectivity between the client workstation and the server. L2TP requires only that the tunnel medium provide packet-oriented point-to-point connectivity. The protocol may be used over media such as ATM, Frame Relay, and X.25. L2TP provides the same functionality as PPTP. Based on Layer 2 Forwarding (L2F) and PPTP specifications, L2TP allows clients to set up tunnels across intervening networks.

- [0039] In another aspect, for different micro-tunnels, the sequence space for the sequence field of the GRE header 204 is different. If all the micro-tunnels share the same sequence space, then the R-P interface may not be able to take advantage of treating the Differentiated Services Code Point (DSCP) marking differently. DSCP is used for implemented Quality of Service (QoS). A replacement header field, called the DS field, includes six bits of as a DSCP codepoint, to select the per-hop-behavior a packet experiences at each node. The DSCP is detailed in RFC 2474, described hereinabove.
- [0040] The receiver would re-order the packets based on the GRE sequence number and any gain which could have been achieved by the R-P interface giving packets with certain code-point a higher priority would be lost. If different micro-tunnels do not share the same sequence space, the PDSN may use a different micro-tunnel for sending packets with different DSCP.
- [0041] Service differentiation for the traffic carried by each micro-tunnel is independent of the micro-tunnel ID and is based on the outer DSCP or other signaling information exchanged between the PDSN and RAN (e.g., RSVP).

Format of the GRE key field:

- [0042] FIG. 3 illustrates the GRE key field 300 of the GRE header 204 according to one embodiment, wherein the GRE key field 300 includes two fields: Packet Service Identifier (PSI) 302; and Micro-Tunnel Identifier (MTID) 304. The boundary 306 between the two fields is not fixed, and therefore is illustrated to indicate the boundary may be adjusted or determined by the PCF/BSC 104 or the PDSN 102. The GRE key field 300 is used by the PDSN 102 to identify the micro-tunnel for a given user by the MTID, as well as identifying the associated MN by the PSI.
- [0043] To build the GRE key field 300, the PCF/BSC 104 receives a request for a data service from a MN, such as MN 108. The PCF/BSC 104 requests the establishment of a link for servicing the data service for MN 108. The PCF/BSC 108 sends a GRE key configuration record to the PDSN 104. The GRE key configuration record may be provided in the form <PSI,L>, wherein L indicates the length of the MTID field 304. For example, for if the record is given as <00,2>, the PSI is identified by digital value 00 and the last two bits are available for identifying the MTID. Each value of the GRE key field in the GRE tunnel between the PDSN 102 and the PCF/BSC 104 identifies a micro-tunnel. For the PSI field determination, the PCF/BSC 104 structures a list of

<network address, subnet mask> pairs from which the PCF chooses to associate a mobile node.

[0044] A general scheme for constructing the GRE key field 300 allows the PCF to determine the available PSI values for a given mobile node. In other words, the PCF determines the GRE key configuration record. In one embodiment, the GRE key field 300 is specified as having a fixed number of bits, i.e., a fixed length. For example, the GRE key field 300 may be specified as 32 bits defining a GRE key space as illustrated in FIG. 4. Each value in the GRE key space is identified by four bits. The GRE key field 300 is used to identify both the PSI field 302 and the MTID field 304 as illustrated in FIG. 3. Therefore, if two bits are used to identify the MN, i.e., the destination identifier PSI, there are two bits left to identify the micro-tunnel, i.e., for the micro-tunnel identifier MTID. In this way, the PCF is able to allocate the total available values in the GRE key space to multiple mobile nodes.

[0045] As an example, the PCF may assign the two MSB bits 00 to MN 108. The configuration record would be <00,2> indicating that 2 bits remain for micro-tunnel identifiers, i.e., MTID. The corresponding GRE key values available for MN 108 are then 0000, 0001, 0010, and 0011. The MN 108 would have 4 available identifiers for micro-tunnels. The PCF may then assign the three MSB bits 010 to MN 110, wherein the configuration record would be <010,1> indicating there is one bit left for the MTID. In this case, MN 110 would have 2 identifiers available for micro-tunnels. The resultant GRE key values available for MN 110 would be 0100, 0101. In other words, the boundary 306 between the PSI field 302 and the MTID field 304 is variable per mobile node. The ability to craft the PSI and MTID fields 302, 304 available for different mobile nodes may result in a GRE key space which is fragmented. The fragmentation provides flexibility in resource allocation within the system. As described hereinabove, the PCF determines the assignments within the GRE key space and provides this information to the PDSN in the form of a configuration record.

[0046] It is desirable to allocate the available identifiers for multiple mobile nodes, and therefore, the PCF determines a range of values for each mobile node. Such determination may be based on historical usage of mobile nodes, available services, or some other design criteria specific to the system. While the GRE key field 300 is specified as a fixed length, the PSI and MTID fields 302, 304 have variable length, as indicated by the variable boundary 306. The longer the PSI field 302, i.e., more bits

allocated to PSI, the more mobile nodes may be identified, as the PSI is used to identify the mobile node. The longer PSI fields, however, leave fewer bits for the MTID, which identifies each of the micro-tunnels for a given mobile node, and therefore, the fewer micro-tunnels available per mobile node. Similarly, shorter PSI fields allow fewer MNs, but allow more micro-tunnels per MN.

[0047] Note, alternate embodiments may utilize an alternate field having a different number of bits than the GRE key field 300. Still other embodiments may implement a field having a variable number of bits, the PSI and MTID fields 302, 304 are then allocated within the variable length field. In these latter embodiments, the PSI and MTID length allocation may be determined proportionally, or may be specifically determined given the current length of the variable length field.

[0048] When the PDSN 102 receives traffic destined for the PPP instance associated with a particular mobile node, the PDSN 102 encapsulates the traffic in a GRE tunnel and sets the GRE key field 300 as described herein. The PDSN 102 sets the Most Significant Bits (MSBs) of the GRE key field 300 (i.e., PSI field 302) to the one of the network addresses which the PCF/BSC 104 has advertised in the A11-Registration-Request message for a particular MN, wherein each PPP instance is associated with an IMSI. The length of the network address is determined by the subnet mask associated with the network address used.

[0049] The PDSN 102 sets the LSBs of the GRE key field 300 (i.e., MTID field 304) to identify the micro-tunnel in which the packet should be carried. The numerical value of the LSBs has no significance and is only used to identify a micro-tunnel.

[0050] The PCF/BSC 104 routes packets received via micro-tunnels to the mobile stations by examining the GRE key field 300 of the GRE header 204 and determining the associated mobile station ID based on the "routing tables" advertised to the PDSN in the A11-Registration-Request message. In one case, the PCF/BSC 102 may specify the MSBs of the GRE key field 300 and allow the PDSN 104 to specify the LSBs of the GRE key field 300.

[0051] In order to make the A11-Registration Message backwards compatible, the PCF/BSC 104 may populate the PSI field in the A11-Registration Request with the PSI field which is left-justified and append the length of the PSI field as a new information element to the A11-Registration Request message.

- [0052] An alternative method of specifying the GRE Key associated with micro-tunnels is where the BSC/PCF sends to the PDSN (in an A11-Registration Request message) the entire 32-bits of the GRE Key for the micro-tunnel along with the QoS characteristics of the micro-tunnel to be established.
- [0053] The PDSN 102 is the entity to drop packets if congestion occurs at the RAN 120, which is where a bottleneck may be expected. Note that the PDSN 102 is the entity which may drop a whole IP packet without the need to remove the link layer framing (the BSC gets the packets when HDLC is already applied to them). Also, the PDSN 102 distinguishes a PPP control packet from a PPP frame containing data (again the PCF/BSC 102 has to peek into the packet in order to make this distinction). The RAN is where the queues associated with packets with different QoS requirements are formed.
- [0054] Because of the above facts, the PCF/BSC 104 may be the entity which provides back-pressure to the PDSN 102. The PCF/BSC 104 should apply back-pressure based on the DS code points. The idea is that the length of the PCF/BSC 104 queues for different DSCPs may be different because the PCF/BSC services the bins for different DSCPs differently. More precisely, the PCF should be able to apply back-pressure by specifying <PSI, DSCP, MTID> triplet. The mobile should be able to set any of the PSI, DSCP, or MTID to a wild-card value. For example, a <PSI, *, *> indicates a back-pressure for all the traffic destined for the mobile identified by the PSI.
- [0055] The current A-interface signaling maps each air interface service instance identified by an sr_id to a GRE tunnel identified by <src_ip = PDSN_IP, dest_ip = PCF_IP, GRE_key = PSI>. It is also expected for the PDSN to map the received packets from the Internet side to an appropriate air-interface pipe which is identified by the sr_id. In the method, the following assumptions are made: (1) the air-interface pipes (aka service instances) with appropriate QoS characteristics are assumed to be established prior to reception of the IP packets; (2) the PDSN is assumed to be aware of the QoS characteristics of these service instances.
- [0056] The first assumption is wasteful as there may be no packets received at the PDSN for the microflow. Also, the number of possible configurations of the air-interface streams may be well above the maximum number of simultaneous streams which the air-interface may support. Therefore, it may not be possible or desirable to pre-establish air-interface streams with all the possible configurations and expect the

PDSN to map the flows to the air-interfaces streams. The second assumption requires a new set of signaling to transfer information (i.e., characteristics of the air-interface streams) from the Radio Access Network (RAN) to the PDSN. This may be avoided when the mapping of the micro-flows to air-interface streams is performed at the RAN.

[0057] The RAN has access to two pieces of information used to create air-interface pipes with the required characteristics and route the received packets from the PDSN to these air-interface pipes: a) The RAN knows what type of service instance are to be created based on the QoS characteristics of the packet given by the PDSN; and b) The RAN knows the characteristics of the service instances.

[0058] Another consideration is authorization. The RAN, rather than the PDSN, is typically the entity providing the air-interface QoS as a service to the user. Therefore, it is reasonable to assume the RAN is responsible for authorizing the user to receive QoS as a service or not based on a user's RAN subscription. Additionally, it is desirable for the RAN to be the entity which decides whether a new service instance will be created to differentiate between handling of packets.

[0059] The PDSN receives packets from different subnets on the Internet side. The PDSN may have a different Service Level Agreements (SLA) with these subnets. These SLAs enable the PDSN to understand how each packet is to be treated depending on the DSCP value. The BSC only understand a certain mapping between the DSCP and Per-Hop Behavior (PHB) which is given to it by the PDSN. Therefore, the PDSN should translate (i.e., re-map) the received DSCPs into code-point values which are understood by the BSC. This is a functionality expected from any other edge-router in the DS frame-work.

[0060] In another embodiment the mobile knows the QoS requirements associated with a particular micro-flow. The mobile tells the BSC/RAN the requested QoS requirements and the GRE key associated with this QoS treatment. The RAN then determines whether the mobile is authorized to receive this treatment. The RAN may verify authorization by checking with an Accounting, Authentication, and Authorization (AAA) server to determine the client's authorization based on a subscription. Alternately, the RAN may make this determination based on load. The RAN may then choose to establish an air-interface service instance (e.g., an RLP instance) which is appropriate for the QoS treatment requested (e.g., no RLP re-transmission for delay-sensitive QoS requirement). The RAN then requests establishment of the A10/A11

interfaces which is going to carry this microflow. The MS then conveys to the PDSN the required QoS characteristics and the information which specifies the micro-flow (e.g., the source IP address, destination IP address, protocol, and the port number).

[0061] In yet another embodiment, if the PDSN knows a certain micro-flow requires a particular QoS treatment, that information is conveyed either directly through the IP pipe between PDSN and the mobile, or through the A-interface and then the air-interface signaling. The rest of the scenario then becomes similar to the scenario in which the mobile is the entity aware of the QoS requirements for the micro-flow. The PDSN may become the entity aware of the QoS requirements for a micro-flow based on the SLA as described hereinabove.

[0062] New A-interface messaging is needed for the PDSN to identify to the BSC the characteristics associated with packets carried by certain micro-tunnels. For example, the PDSN indicates to the PCF that packets carried by micro-tunnel <PSI, MTID> are zero-byte header compressed. The BSC uses this information to carry these packets on the air-interface using an "RLP-free" service instance. In particular, the PDSN should be able to associate the following attributes to each micro tunnel:

- Dropping is permitted or not within this micro-tunnel. This is a key Quality of Service criterion.
- Link-Layer Assisted ROHC (LLAROH) (zero-byte header compression) application. This information is needed by the BSC for routing the packets received from the PDSN to the service instance negotiated for zero-byte header compression as explained hereinabove.
- Mapping between DSCPs and the desired Per-Hop Behavior (PHB). In IP packet headers, PHBs are not indicated as such; instead Differentiated DSCP values are used. In a given network domain, there is a locally defined mapping between DSCP values and PHBs. Standardized PHBs recommend a DSCP mapping, but network operators may choose alternative mappings.
- Signaling messages allowing the PDSN to ask if the RAN supports traffic which requires a certain type of treatment. Such query may include RSVP messages which enable the PDSN to perform admission control for specific traffic.

- [0063] Those of skill in the art would understand that information and signals may be represented using any of a variety of different technologies and techniques. For example, data, instructions, commands, information, signals, bits, symbols, and chips that may be referenced throughout the above description may be represented by voltages, currents, electromagnetic waves, magnetic fields or particles, optical fields or particles, or any combination thereof.
- [0064] Those of skill would further appreciate that the various illustrative logical blocks, modules, circuits, and algorithm steps described in connection with the embodiments disclosed herein may be implemented as electronic hardware, computer software, or combinations of both. To clearly illustrate this interchangeability of hardware and software, various illustrative components, blocks, modules, circuits, and steps have been described above generally in terms of their functionality. Whether such functionality is implemented as hardware or software depends upon the particular application and design constraints imposed on the overall system. Skilled artisans may implement the described functionality in varying ways for each particular application, but such implementation decisions should not be interpreted as causing a departure from the scope of the present invention.
- [0065] The various illustrative logical blocks, modules, and circuits described in connection with the embodiments disclosed herein may be implemented or performed with a general purpose processor, a digital signal processor (DSP), an application specific integrated circuit (ASIC), a field programmable gate array (FPGA) or other programmable logic device, discrete gate or transistor logic, discrete hardware components, or any combination thereof designed to perform the functions described herein. A general purpose processor may be a microprocessor, but in the alternative, the processor may be any conventional processor, controller, microcontroller, or state machine. A processor may also be implemented as a combination of computing devices, e.g., a combination of a DSP and a microprocessor, a plurality of microprocessors, one or more microprocessors in conjunction with a DSP core, or any other such configuration.
- [0066] The steps of a method or algorithm described in connection with the embodiments disclosed herein may be embodied directly in hardware, in a software module executed by a processor, or in a combination of the two. A software module may reside in RAM memory, flash memory, ROM memory, EPROM memory,

EEPROM memory, registers, hard disk, a removable disk, a CD-ROM, or any other form of storage medium known in the art. An exemplary storage medium is coupled to the processor such the processor may read information from, and write information to, the storage medium. In the alternative, the storage medium may be integral to the processor. The processor and the storage medium may reside in an ASIC. The ASIC may reside in a user terminal. In the alternative, the processor and the storage medium may reside as discrete components in a user terminal.

[0067] The previous description of the disclosed embodiments is provided to enable any person skilled in the art to make or use the present invention. Various modifications to these embodiments will be readily apparent to those skilled in the art, and the generic principles defined herein may be applied to other embodiments without departing from the spirit or scope of the invention. Thus, the present invention is not intended to be limited to the embodiments shown herein but is to be accorded the widest scope consistent with the principles and novel features disclosed herein.

WHAT IS CLAIMED IS: